# ProbLog: A Probabilistic Prolog and its Application in Link Discovery 

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## Paper context

- At the time when the paper was released, there were no programs for modelling the exact inference for discrete variables
- Discrete variables require separate rules than the continuous variables


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## Intro to Prolog

- Part of the logical programming languages family
- Program consists of a set of definite clauses
- Programs can contain the following: rules, facts and variables
- Clauses can be only True or False


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## Prolog example

```
burglary.
hears_alarm(mary).
alarm :- burglary.
alarm :- earthquake.
calls(X) :- alarm, hears_alarm(X).
call :- calls(X).
9
```

- Alarm and calls are called rules
- hears_alarm, burglary are called facts
- Mary is a variable


## Why extending Prolog to Probabilistic Programming?

- Adding probabilities to clauses is closer to real-world problems
- Probabilistic Database is slow -> 10 or more conjuncts are infeasible to compute
- Many practical applications (i.e. life sciences) require computing probabilities in network relations


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## Intro to ProbLog

- Built on top of Prolog, both being very similar
- Only major difference: Problog has probabilities of success attached to the clauses
- It has equivalent functions for sample and observe (can you spot them in the next slide?)


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## ProbLog example

```
    1.0:: likes(X,Y):- friendof( }X,Y)\mathrm{ .
0.8:: likes(X,Y):- friendof(X,Z), likes(Z,Y).
0.5:: friendof(john,mary).
0.5:: friendof(mary, pedro).
0.5:: friendof(pedro,tom).
What are sample and observe here?
evidence(likes(john, pedro), false).
query(likes(mary, tom)).|
\begin{tabular}{lll} 
Query \(\boldsymbol{V}\) & Location & Probability \\
\hline likes(mary,tom) & \(11: 7\) & 0.15 \\
\hline
\end{tabular}
```

Screenshots taken from:
https://dtai.cs.kuleuven.be/problog/tutorial/basic/02_b ayes.html (more examples there as well)

## Computing queries

Two steps:

1. Build monotone DNF formula representing all solutions
2. Compute the probability of this DNF formula

## Computing queries

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1. Build monotone DNF formula representing all solutions

- SLD-resolution to transform query into equivalent tree

Root is query to be proven
Recursively generate subgoals

- Use the disjunction of proof paths in tree as DNF

2. Compute the probability of this DNF formula

## Computing queries | SLD-resolution example

```
        ?- \(1(\mathrm{j}, \mathrm{t})\).
        ll/ \(\quad 12\)
\(:-\mathrm{fo}(\mathrm{j}, \mathrm{t}) .:-\mathrm{fo}(\mathrm{j}, \mathrm{A}), \mathrm{l}(\mathrm{A}, \mathrm{t})\).
            fl 1
            :- \(1(\mathrm{~m}, \mathrm{t})\).
\(:-\frac{l \nu}{l l} \quad:-\mathrm{fo}(\mathrm{m}, \mathrm{B}), \mathrm{l}(\mathrm{B}, \mathrm{t})\).
```



```
\(:-\mathrm{fo}(\mathrm{p}, \mathrm{t}) . \quad:-\mathrm{fo}(\mathrm{p}, \mathrm{C}), \mathrm{l}(\mathrm{C}, \mathrm{t}) . \quad:-\mathrm{fo}(\mathrm{t}, \mathrm{t}) . \quad: \quad:-\mathrm{fo}(\mathrm{t}, \mathrm{E}), \mathrm{l}(\mathrm{E}, \mathrm{t})\).
\(f 41 \quad f 41\)
            \(\square \quad l l)^{:-1(t, t)} l_{2}\)
                \(: \underline{i-\mathrm{fo}(\mathrm{t}, \mathrm{t}) .} \quad: \underline{-\mathrm{fo}(\mathrm{t}, \mathrm{D}), \mathrm{l}(\mathrm{D}, \mathrm{t}) .}\)
```

1.0: likes $(X, Y)$ :- friendof $(X, Y)$.
0.8: likes $(X, Y)$ :- friendof $(X, Z)$, likes $(Z, Y)$.
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## Computing queries | SLD-resolution example


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$$
P\left(\left(l_{1} \wedge l_{2} \wedge f_{1} \wedge f_{2} \wedge f_{1}\right) \vee\left(\underline{\left.l_{1} \wedge l_{2} \wedge f_{1} \wedge f_{3}\right)}\right) .\right.
$$

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## Computing queries

Two steps:

1. Build monotone DNF formula representing all solutions
2. Compute the probability of this DNF formula

- Using Binary Decision Diagram (BDD) representation

Start from full binary tree, merging isomorphic subgraphs and deleting redundant nodes

## Computing queries | BDD calculation example



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$\begin{array}{llllll}0 & 1 & 2 & 3 & 4 & 5\end{array}$

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$$
P\left(\left(l_{1} \wedge l_{2} \wedge f_{1} \wedge f_{2} \wedge f_{4}\right) \vee\left(l_{1} \wedge l_{2} \wedge f_{1} \wedge f_{3}\right)\right)
$$

## Computing queries | BDD calculation example

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0
2
3
4
5
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$$
P\left(\left(l_{1} \wedge l_{2} \wedge f_{1} \wedge f_{2} \wedge f_{4}\right) \vee\left(l_{1} \wedge l_{2} \wedge f_{1} \wedge f_{3}\right)\right)
$$

## Computing queries | BDD calculation example

$$
P\left(\left(l_{2} \wedge f_{1} \wedge f_{2} \wedge f_{4}\right) \vee\left(l_{2} \wedge f_{1} \wedge f_{3}\right)\right)
$$

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## Computing queries

Two steps:

1. Build monotone DNF formula representing all solutions
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Start from full binary tree, merging isomorphic subgraphs and deleting redundant nodes

- Heuristically determine variable order in SOTA BDD algorithms
- Reusable BDD for different queries


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## Approximating the success probability

- Why approximate?
- Iterative deepening to compute SLD-tree
- Use incomplete SLD-tree to derive upper and lower bound
- Lower bound encodes successful proofs found so far
- Upper bound encodes all proofs all proofs found so far
- Keep growing tree until upper and lower bound are sufficiently close


## Results



- Good runtime in terms of level depthness
- Can deal with many conjuncts, up to 100k.
- Probability is converging to the true one after the 6th depth level
- Bounds are converging to $\sim 0.2$ after the 6th depth level.

Running times for 10 test graphs with 1400 edges.

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## Results



Convergence of the probability interval for 10 test graphs with 1400 edges.


Convergence of bounds for one graph with 1800 edges, as a function of the search level.

## Questions

- What is the addition of ProbLog to Prolog?


## Questions

- What is the addition of ProbLog to Prolog?
- What other probabilistic programming languages also have inherently included upper and lower bounds?


## Thank you for your attention!

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